Category-Theoretic Reconstruction of Schemes from Categories of Reduced Schemes

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Today's Talk

Let S be a scheme, Φ/S a set of properties of S-schemes, and $\mathrm{Sch}_{\Phi/S}$ the full subcategory of $\mathrm{Sch}_{/S}$ determined by the objects $X \in \mathrm{Sch}_{\Phi/S}$ that satisfy every property of Φ/S .

In this talk, I will explain how to reconstruct S from $Sch_{\Phi/S}$.

Notations and Conventions

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S: \mathsf{Scheme} \blacklozenge/S: \mathsf{a} \ \mathsf{set} \ \mathsf{of} \ \mathsf{properties} \ \mathsf{of} \ S\mathsf{-schemes} \mathsf{Sch}_{\blacklozenge/S}: \begin{cases} \mathsf{the} \ \mathsf{full} \ \mathsf{subcategory} \ \mathsf{of} \ \mathsf{Sch}_{/S} \ \mathsf{determined} \ \mathsf{by} \\ \mathsf{the} \ \mathsf{objects} \ X \in \mathsf{Sch}_{\blacklozenge/S} \ \mathsf{that} \ \mathsf{satisfy} \ \mathsf{every} \ \mathsf{property} \ \mathsf{of} \ \blacklozenge/S \\ \times, \lim: \mathsf{the} \ \mathsf{fiber} \ \mathsf{product}, \ \mathsf{limit} \ \mathsf{in} \ \mathsf{Sch} \\ \times^{\blacklozenge}, \lim^{\blacklozenge}: \mathsf{the} \ \mathsf{fiber} \ \mathsf{product}, \ \mathsf{limit} \ \mathsf{in} \ \mathsf{Sch}_{\blacklozenge/S} \end{cases}
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In the present talk, we shall mainly be concerned with the properties

$$\blacklozenge \subset \{\mathrm{red}, \mathrm{qcpt}, \mathrm{qsep}, \mathrm{sep}\}.$$

Previous Research

Mochizuki 2004 : $\phi/S = \text{f.t.}/S$, S: locally Noetherian

(+ log scheme version)

van Dobben de Bruyn 2019 : $\spadesuit = \varnothing$, S: arbitrary

Wakabayashi 2010 : superscheme version of the case of Mochizuki

Anabelian Geometry $: \phi/S = \text{f\'et}/S$

These research and my research are motivated by anabelian geometry.

Main Theorem

Main Theorem

- (1) S: locally Noetherian normal scheme, $\blacklozenge \subset \{\text{red}, \text{qcpt}, \text{qsep}, \text{sep}\}$. Then the following may be reconstructed category-theoretically from $\mathsf{Sch}_{\blacklozenge/S}$:
 - (a) the structure of T as a scheme (for every object $T \in \operatorname{Sch}_{\blacklozenge/S}$),
 - (b) the structure of f as a morphism of schemes (for every morphism $(f:X\to Y)\in \operatorname{Sch}_{\Phi/S}$).
- (2) S, T: quasi-separated,
 - $\blacklozenge, \lozenge \subset \{\text{red}, \text{qcpt}, \text{qsep}, \text{sep}\} \not\subset \blacklozenge, \{\text{qsep}, \text{sep}\} \not\subset \lozenge$ Then, $\mathsf{Sch}_{\blacklozenge/S} \cong \mathsf{Sch}_{\lozenge/T} \Rightarrow \blacklozenge = \lozenge$.
- (3) S,T: locally Noetherian normal schemes, $\blacklozenge \subset \{\mathrm{red}, \mathrm{qcpt}, \mathrm{qsep}, \mathrm{sep}\}$. Then, the following natural functor is equivalent:

$$\operatorname{Isom}(S,T) \to \operatorname{\mathbf{Isom}}(\operatorname{\mathsf{Sch}}_{\blacklozenge/T},\operatorname{\mathsf{Sch}}_{\blacklozenge/S})$$
$$f \mapsto f^*$$

Outline

Since a scheme is constructed by

- the underlying set,
- the underlying topological space, and
- the structure sheaf,

to reconstruct a scheme, it suffices to reconstruct these structures.

In the present talk, I explain how to reconstruct the underlying sets, and give category-theoretic characterizations of various properties used to reconstruct the underlying topological spaces and the structure sheaves.

Remark: the Fiber Product in $Sch_{\bullet/S}$

Lemma

 $f:Y \to X, g:Z \to X$: morphisms in $Sch_{\P/S}$. Suppose that either f or g is quasi-compact.

Then, the fiber product $Y\times_X^{\blacklozenge}Z$ in $\mathrm{Sch}_{\blacklozenge/S}$ exists, and the following assertions hold:

If $\operatorname{red} \not\in \spadesuit$, then $Y \times_X^{\spadesuit} Z \cong Y \times_X Z$.

If $red \in \blacklozenge$, then $Y \times_X^{\blacklozenge} Z \cong (Y \times_X Z)_{red}$.

In particular, $Y \times_X Z$ and $Y \times_X^{\blacklozenge} Z$ have same underlying top.

An Idea to Reconstruct the Underlying Sets

Observation

A point $x \in X$ may be determined by

$$f: Y \to X$$
 s.t. $|Y|$: 1pt. set, and $Im(f) = \{x\}$.

Hence,

giving a point of $X \iff$

giving a certein equivalence class of $f:Y\to X$ s.t. |Y|: 1pt. set.

To reconstruct the underlying set, it suffices to characterize one-pointed schemes (i.e., schemes whose underlying sets are 1pt. sets) cat.-theoretically.

Characterization of the One-Pointed Schemes

Let $X \in \mathsf{Sch}_{\blacklozenge/S}$.

Characterization of the 1pt. Scheme

|X| is **not** 1pt. set \iff

$$\exists Y,Z\neq\varnothing\;,\;\exists Y\to X,Z\to X\quad\text{s.t.}\quad Y\times_X^{\blacklozenge}Z=\varnothing$$

..)

X has two distinct pts. $x_1, x_2 \Rightarrow \operatorname{Spec}(k(x_1)) \times_X^{\blacklozenge} \operatorname{Spec}(k(x_2)) = \emptyset$.

X satisfies the condition $\Rightarrow y \in Y, z \in Z$ determine two distinct pts. of X.

Reconstruction of the Underlying Set 1

Let $X \in Sch_{\phi/S}$. We define

$$\operatorname{Pt}_{\Phi/S}(X) := \left\{ (p_Z : Z \to X) \in \operatorname{Sch}_{\Phi/S} \mid |Z| : \text{ 1pt. set} \right\} / \sim,$$

where

$$(p_Z:Z\to X)\sim (p_{Z'}:Z'\to X):\stackrel{\mathsf{def}}{\iff} Z\times_{p_Z,X,p_{Z'}}^{\bullet}Z'\neq\varnothing.$$

Reconstruction of the Underlying Set

 $\operatorname{Pt}_{lacklet/S}:\operatorname{Sch}_{lacklet/S} o\operatorname{Set}$ is naturally isomorphic to the functor $U_{lacklet/S}^{\operatorname{Set}}:\operatorname{Sch}_{lacklet/S} o\operatorname{Set}.$

Reconstruction of the Underlying Set 2

Since the functor $\mathsf{Pt}_{\blacklozenge/S}$ is defined category-theoretically, the following corollary holds:

Corollary

If $F: \mathsf{Sch}_{\Phi/S} \to \mathsf{Sch}_{\Diamond/T}$ is an equivalence, then $U_{\Phi/S}^{\mathsf{Set}} \cong U_{\Diamond/T}^{\mathsf{Set}} \circ F$.

Regular Monomorphisms

 \mathcal{C} : category, $(f: X \to Y) \in \mathcal{C}$.

Definition

f is a **regular monomorphism**

 $: \stackrel{\mathsf{def}}{\Longleftrightarrow} \ \exists g,h: Y \to Z, \text{ s.t., } f \text{ is the equalizer of } (g,h).$

Property of reg. mono. in $Sch_{\blacklozenge/S}$

 $S \colon \operatorname{q.s.}, \ (f:X \to Y) \in \operatorname{Sch}_{\blacklozenge/S} \colon \operatorname{reg. mono.} \ \Rightarrow f \colon \operatorname{immersion}.$

 \therefore) f: reg. mono. \Rightarrow f: b.c. of the diagonal (details omitted).

Corollary (Cat.-Theoretic Characterization of Red. Schemes)

 $X \in \mathsf{Sch}_{\Phi/S}$ is red. $\iff [f: Y \to X: \mathsf{surj. reg. mono.} \Rightarrow f: \mathsf{isom.}]$

:) a surj. reg. mono. is a surj. closed immersion.

Closed Immersions

Closed immersions may be characterized as follows:

Characterization of Closed Immersions

S: q.s., $(f: X \to Y) \in \operatorname{\mathsf{Sch}}_{\blacklozenge/S}$. f: closed immersion if and only if

- f: reg. mono.
- $\forall (T \to Y)$, the b.c. $X_{\blacklozenge,T} = X \times_Y^{\blacklozenge} T$ exists.
- $\forall (T \to Y)$, $\forall t \in T$: closed pt. s.t. $t \notin \text{Im}(f_{\blacklozenge,T} : X_{\blacklozenge,T} \to T)$, $X_{\blacklozenge,T} \coprod \text{Spec}(k(t)) \to T$: reg. mono.

Hence to give a cat.-theoretic characterization of closed immersions, it suffices to characterize the closed pt.

In particular, it suffices to characterize the relation $x_1 \rightsquigarrow x_2$.

Strongly Local 1

S: q.s., $X \in \mathsf{Sch}_{\blacklozenge/S}$, $x_1, x_2 \in X$.

Definition (Strongly Local)

 (X, x_1, x_2) is strongly local in $Sch_{\phi/S} : \stackrel{\mathsf{def}}{\Longleftrightarrow}$

- X: connected
- $\forall (f:Z\to X)$: reg. mono., $[x_1,x_2\in \mathrm{Im}(f),\Rightarrow f$: isom.].
- $\operatorname{Spec}(k(x_1)) \coprod \operatorname{Spec}(k(x_2)) \to X$: epi.
- Spec $(k(x_1)) \to X$: reg. mono.
- $\forall (f:Z\to X)$: reg. mono., $[x_1\not\in \mathrm{Im}(f), Z\neq\varnothing\Rightarrow Z\coprod \mathrm{Spec}(k(x_1))\to X$: **not** a reg. mono.].

Remark

The property that (X, x_1, x_2) is strongly local is defined cat.-theoretically from the data $(\operatorname{Sch}_{\Phi/S}, X, x_1, x_2)$.

Strongly Local 2

 $S: \mathsf{q.s.}, \ X \in \mathsf{Sch}_{\blacklozenge/S}, x_1, x_2 \in X.$

Properties of Strongly Local Objects

If (X, x_1, x_2) : strongly local, then

- (1) $X \cong \operatorname{Spec}(\operatorname{local domain})$
- (2) One of x_1, x_2 is the closed pt., and the other is the generic pt. In particular, $x_1 \rightsquigarrow x_2$ or $x_2 \rightsquigarrow x_1$.

Let $V = \operatorname{Spec}(\text{valuation ring}), v \in V$: closed pt., $\eta \in V$: generic pt.

Proposition (Spec. of Valuation Rings are Strongly Local)

 (V, v, η) : strongly local.

Cat.-Theoretic Characterization of " $x_1 \rightsquigarrow x_2$ or $x_2 \rightsquigarrow x_1$ "

 $S: \text{ q.s., } X \in \mathsf{Sch}_{\blacklozenge/S}, \ x_1, x_2 \in X.$

Cat.-Theoretic Characterization of " $x_1 \rightsquigarrow x_2$ or $x_2 \rightsquigarrow x_1$ ".

" $x_1 \leadsto x_2 \text{ or } x_2 \leadsto x_1$ " \iff

 $\exists Z \in \mathsf{Sch}_{\blacklozenge/S}, \exists z_1, z_2 \in Z, \exists (f:Z \to X) \in \mathsf{Sch}_{\blacklozenge/S}, \text{ s.t., } \\ (Z,z_1,z_2): \text{ str. loc., and } \{f(z_1),f(z_2)\} = \{x_1,x_2\}.$

By using the above characterization, we can characterize the relation $x_1 \rightsquigarrow x_2$ (details omitted).

Corollary

- (1) Closed immersions may be characterized cat.-theoretically.
- (2) Underlying top. may be reconstructed cat.-theoretically.

In particular, top.-theoretic properties of schemes (or morphisms) may be characterized cat.-theoretically (ex: q.s., q.c., sep., irred., local ($\cong \operatorname{Spec}(\operatorname{local\ ring})$), open imm., univ. closed, etc.). 16 / 30

Reconstruction of the Underlying Top.

(Similarly to the case of Set) $\forall F: \mathsf{Sch}_{\blacklozenge/S} \xrightarrow{\sim} \mathsf{Sch}_{\lozenge/T}$, the following diagram commutes (up to isom.):

An Observation

To reconstruct the structure sheaf of $X \in \mathsf{Sch}_{\blacklozenge/S}$, it suffices to characterize the ring scheme $\mathbb{A}^1_X \to X$ cat.-theoretically.

Since \mathbb{A}^1 is f.p. over a base scheme, we want to get a cat.-theoretic characterization of f.p. morphisms.

Idea

$$\mathrm{f.p.}/S \rightleftharpoons \mathrm{a}$$
 "compact object" in $\mathrm{Sch}_{/S}^{\mathrm{op}}$

More precisely,

$$X \to S$$
: f.p. \iff

$$\forall (V_{\lambda}, f_{\lambda\mu})_{\lambda \in \Lambda}$$
: diagram in Sch_{/S} s.t.

$$\Lambda$$
: cofiltered, V_{λ} : affine,

the following natural map is surj. :

$$\varphi: \operatorname{colim}_{\lambda \in \Lambda^{\operatorname{op}}} \operatorname{Hom}_{\operatorname{\mathsf{Sch}}_{/S}}(V_{\lambda}, X) \to \operatorname{Hom}_{\operatorname{\mathsf{Sch}}_{/S}}(\lim_{\lambda \in \Lambda}^{\blacklozenge} V_{\lambda}, X).$$

Locally of Finite Presentation Morphisms 1

$$S$$
: q.s., $(f: X \to Y) \in \operatorname{Sch}_{\Phi/S}$, $x \in X$.

Proposition

 $f_x^\#: \mathcal{O}_{Y,f(x)} \to \mathcal{O}_{X,x}$: essentially of finite presentation $\iff \forall (V_\lambda, f_{\lambda\mu})_{\lambda \in \Lambda}$: diagram in $\mathsf{Sch}_{\blacklozenge/Y}$ s.t.

 Λ : cofiltered, V_{λ} : local, $f_{\lambda\mu}(\text{closed pt.}) = f(x)$, the following natural map is surjective :

$$\varphi: \operatornamewithlimits{colim}_{\lambda \in \Lambda^{\operatorname{op}}} \operatorname{Hom}_{\operatorname{\mathsf{Sch}}_{\Phi/Y}}(V_{\lambda}, X) \to \operatorname{Hom}_{\operatorname{\mathsf{Sch}}_{\Phi/Y}}(\varprojlim_{\lambda \in \Lambda}^{\Phi} V_{\lambda}, X).$$

 \therefore) f.p. schemes (over Y) are cpt. objects in Sch_{/Y} (details omitted).

Locally of Finite Presentation Morphisms 2

 $S: \text{ q.s., } (f:X\to Y)\in \mathsf{Sch}_{\blacklozenge/S}.$

Cat.-Theoretic Characterization of Loc.F.P. Morphisms

f: loc. of f.p. \iff

- $\forall x \in X$, $f_x^\#: \mathcal{O}_{Y,f(x)} \to \mathcal{O}_{X,x}$: essentially of finite presentation.
- $\forall (Z \to Y), \forall z \in Z$, the following natural map is bijective :

$$\varphi_{z,X}: \operatorname*{colim}_{W \in I_Z(z)^{\operatorname{op}}} \operatorname{Hom}_{\operatorname{Sch}_{\Phi/Y}}(W,X) \to \operatorname{Hom}_{\operatorname{Sch}_{\Phi/Y}}(\lim_{W \in I_Z(z)}^{\Phi} W,X),$$

where $I_Z(z) := \{i_W : W \to Z \mid i_W : \text{ open imm., } z \in \text{Im}(i_W)\}.$

 \therefore) f.p. schemes (over Y) are cpt. objects in Sch_{/Y} (details composited).

List of Cat.-Theoretic Properties

S: q.s.

 $\forall X \in \mathsf{Sch}_{\Phi/S}, |X|$ has been reconstructed cat.-theoretically, and

the following scheme-theoretic properties have been characterized cat.-theoretically:

- red., irred., integral, q.c., $\cong \operatorname{Spec}(\operatorname{local\ ring})$, $\cong \operatorname{Spec}(\operatorname{field})$.
- q.c., q.s., sep., imm., closed imm., open imm., loc. of f.p., f.p., f.p. + proper (= sep.+ f.p.+ univ. closed).

The following properties have not given yet cat.-theoretic characterizations:

flat, smooth, étale, etc.

An Idea to Reconstruct the Structure Sheaves

To reconstruct the structure sheaf of $X \in \operatorname{Sch}_{\P/S}$, it suffices to characterize the ring scheme $\mathbb{A}^1_X \to X$ cat.-theoretically. Since $\mathbb{A}^1_X = \mathbb{P}^1_X \setminus \{\infty\}$, it suffices to characterize $\mathbb{P}^1_X \to X$ cat.-theoretically.

What to Do

Give a cat.-theoretic characterization of \mathbb{P}^1 .

The Case where $X = \operatorname{Spec}(k)$

$$\mathbb{P}^1_k \iff \begin{cases} \bullet \text{ proper over } \operatorname{Spec}(k) \\ \bullet \text{ the residue field of the generic pt. } \cong k(t) \\ \bullet \text{ "Closest" to } \operatorname{Spec}(k(t)) \end{cases}$$

 \therefore it suffices to characterize $\operatorname{Spec}(k(t)) \to \operatorname{Spec}(k)$. Idea: Lüroth's theorem.

Cat.-Theoretic Characterization of k(t)/k

 $f: Y \to \operatorname{Spec}(k)$: isom. to $\operatorname{Spec}(k(t)) \to \operatorname{Spec}(k)$ over $\operatorname{Spec}(k) \iff$

- $\exists K : \mathsf{field} \ , \ Y \cong \mathrm{Spec}(K)$
- f: not f.p. ($\Leftrightarrow K/k$: not a finite extension)
- $k \subsetneq \forall L \subset K$, \exists isom. $K \cong L$ over k (Lüroth's theorem).
- \leadsto We obtain a cat.-theoretic characterization of $\mathbb{P}^1_k \to \operatorname{Spec}(k)$.

An Idea in the Case of General Base Scheme 1

To characterize $\mathbb{P}^1_X \to X$, it suffices to characterize $\mathbb{P}^1_S \in \operatorname{Sch}_{\Phi/S}$. Since

$$\mathbb{P}^1_S \iff \mathbb{P}^1\text{-bundle}/S + \exists \text{ 3 sections } s_1, s_2, s_3 \text{ s.t. } s_i \cap s_j = \varnothing, (i \neq j)\text{,}$$

it suffices to characterize the \mathbb{P}^1 -bundle over S.

 $\rightsquigarrow \mathbb{P}^1$ -bundle \Rightarrow each fiber is \mathbb{P}^1 .

Remark

- If $red \in \Phi$, then cat.-theoretic fiber \ncong scheme-theoretic fiber.
- A generic fiber may be presented by a limit of open immersions.
 - \leadsto cat.-theoretic generic fiber \cong scheme-theoretic generic fiber.

An Idea in the Case of General Base Scheme 2

 $\mathbb{A}^1 = \mathbb{P}^1 \setminus \{\infty\}$ has a ring scheme structure:

Observation

1-dim ring scheme $= \mathbb{A}^1$??

Lemma (♠♠♠)

R: DVR, $V := \operatorname{Spec} R$, $K := \operatorname{Frac}(R)$ $f : X \to V$: flat ring scheme /V. If f satisfies the following conditions, then $X \cong \mathbb{A}^1_V$ and f is the proj.:

- ullet The special fiber of f is connected and 1-dim.
- The generic fiber of f is \mathbb{A}^1_K .

Without connectedness of the special fiber, there is a counterexample: $\operatorname{Spec}(R[x,(x^{p^2}-x^p)/\pi]).$

The Case of General Base Scheme

Theorem

- S: locally Noetherian normal, $\lozenge = \blacklozenge \cup \{\mathrm{red}\}$, $(f: X \to S) \in \mathsf{Sch}_{\blacklozenge/S}$. f is isom. to $\mathbb{P}^1_S \to S \iff f$ satisfies the following conditions:
- (1) f: f.p. proper.
- (2) $\forall s \in S, f^{-1}(s)_{\text{red}} \cong \mathbb{P}^1_{k(s)}$.
- (3) $\forall \text{generic pt. } \eta \in S, \ f^{-1}(\eta) \cong \mathbb{P}^1_{k(\eta)}.$
- (4) $\exists s_0, s_1, s_\infty$: sections of f s.t. $s_i \cap s_j = \emptyset, (i \neq j)$.
- (5) $\forall i=0,1,\infty$, \exists a ring structure on $X\setminus s_i$ over S in $\mathsf{Sch}_{\lozenge/S}$ s.t. s_j : add. unit, s_k : mult. unit, and $\{i,j,k\}=\{0,1,\infty\}$.
- (6) $(g:Y\to S)\in \operatorname{Sch}_{\Phi/S},\ t_0,t_1,t_\infty$: sections, s.t. satisfy (1),...,(5), $\Rightarrow \exists !h:X\to Y$: closed imm. s.t. $\forall i=0,1,\infty,f=g\circ h,h\circ s_i=t_i.$

Proof

If \mathbb{P}^1 satisfies (6), then by the uniqueness of (6), " \Leftarrow ": ok.

 \therefore It suffices to prove " \Rightarrow " (i.e., \mathbb{P}^1_S satisfies (6)).

Y: satisfies (1),...,(5). We define

$$\begin{split} C: \mathsf{Sch}^{\mathrm{op}}_{/S} \to \mathsf{Set}, \\ (T \to S) \mapsto \left\{ \left. i : \mathbb{P}^1_T \to Y_T \;\middle|\; \begin{array}{c} i : \; \mathsf{closed \; imm.,} \\ 0, 1, \infty \mapsto t_0, t_1, t_\infty \end{array} \right\}. \end{split}$$

Then,

- C: algebraic space /S.
- by (2), each fiber of $C \to S$ is a 1-pt. set.
- by (3), $C \to S$ is birational.

What to Prove

C(S): 1-pt. set.

Proof

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W : \stackrel{\mathsf{def}}{=} \operatorname{Spec}(\mathsf{DVR}).
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\begin{array}{l} \forall (W \to S), \ (Y_W)_{\mathrm{red}} \setminus t_{i,W} \text{: flat ring scheme }/W. \\ \therefore \ \forall W, \ C(W) \text{: 1-pt. set } (\Rightarrow C \text{ is a scheme}). \\ \text{By lemma } (\spadesuit \spadesuit \spadesuit) \text{ and a valuative criterion,} \\ C \to S \text{: proper birat. bij. } (\Rightarrow \text{finite}). \\ \text{Since } S \text{ is normal, } C_{\mathrm{red}} \cong S \text{ (by ZMT).} \end{array}
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In particular, C(S): 1-pt. set. \rightsquigarrow Q.E.D.

The Main Result

(Similarly to the case of Set, Top) $\forall F: \mathsf{Sch}_{\blacklozenge/S} \xrightarrow{\sim} \mathsf{Sch}_{\lozenge/T}: \mathsf{equiv.},$ the following diagram commutes (up to isom.):

Moreover, the following equiv. holds:

$$\operatorname{Isom}(S,T) \xrightarrow{\sim} \operatorname{\mathbf{Isom}}(\operatorname{\mathsf{Sch}}_{\blacklozenge/T},\operatorname{\mathsf{Sch}}_{\blacklozenge/S}).$$

Related Works

I also confirmed that the following problem has been solved:

ullet Reconstructing a Noetherian scheme S from the category of finite S-schemes.

Since we may consider many properties of schemes, there are many cat.-theoretic reconstruction problems.

Thank you for your attention.